

INCREMENTAL DETERMINISTIC PUBLIC- KEY ENCRYPTION

Ilya Mironov, Omkant Pandey,
Omer Reingold, Gil Segev

Microsoft Research

Incremental Deterministic Public-Key Encryption

Deterministic Public-Key Encryption

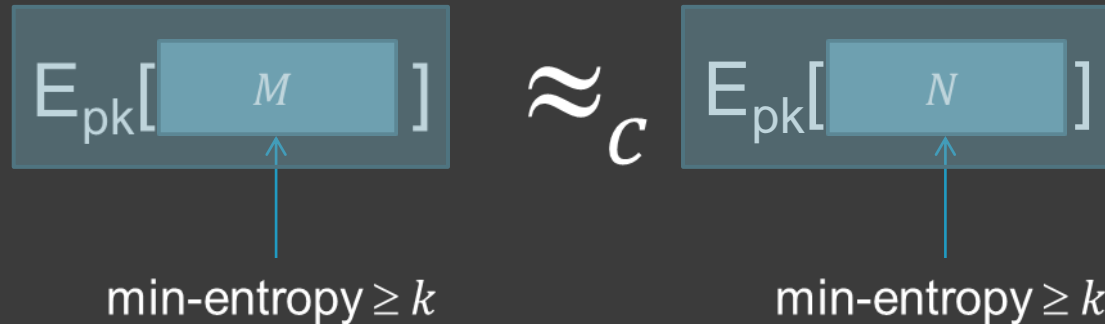
k -source adversary



min-entropy $\geq k$:

Probability of any output $\leq 2^{-k}$

Deterministic Public-Key Encryption: PRIV1-IND



M and N are **independent** of PK

Is It Secure?

- search
- de-duplication
- deterministic KEM

Secure Deterministic Encryption

Computational
assumptions

Min-entropy of
the source

Long, unpredictable plaintext:

- digital photograph
- MS Word document
- entire database
- full disk

security

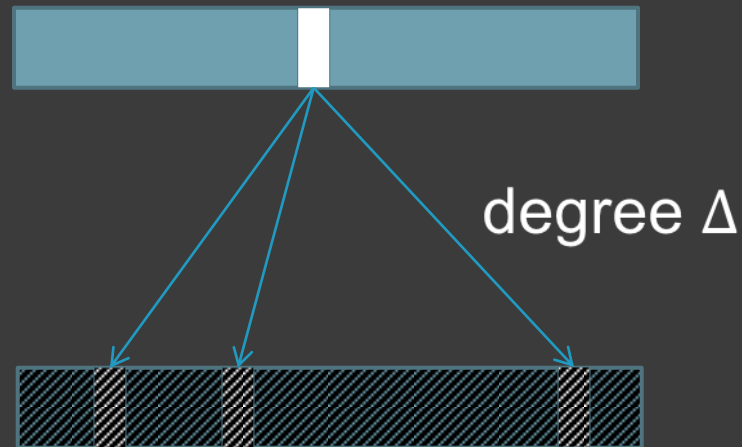


Length of the plaintext

efficiency



Incrementality



- Incrementality with access to plaintext: setting bit
- Incrementality without access to plaintext: flipping bit

Incremental Deterministic Public-Key Encryption

Our results

- Lower bound:

incrementality

$$\Delta > \frac{|P|}{k \log |C|}$$

min-entropy

- Two schemes

- Generic Solution

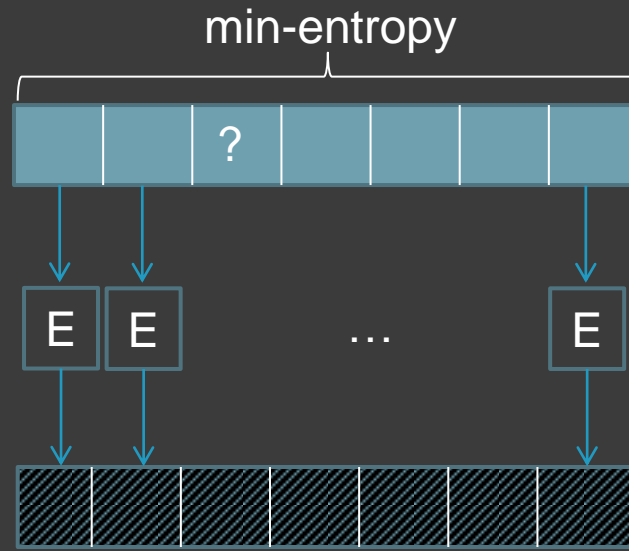
Deterministic Encryption



Incremental Deterministic Encryption

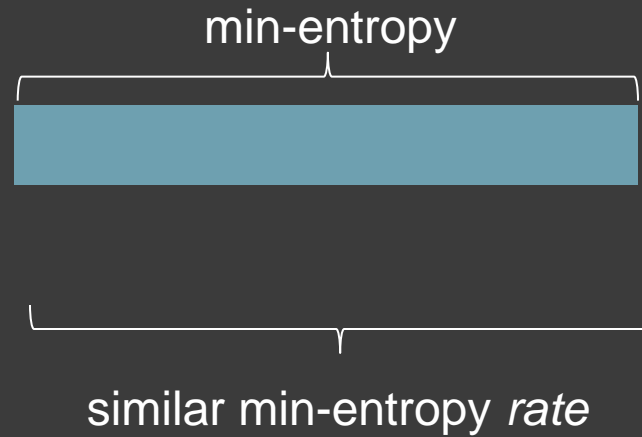
- DDH-based solution
tight up to polylog factors

Naïve Generic Solution



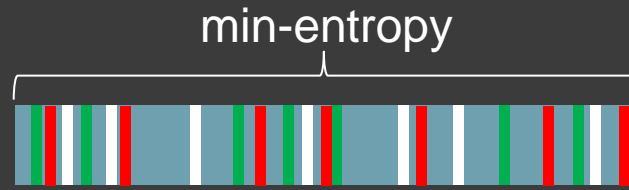
E: deterministic encryption scheme

Sample-then-extract



Generic Solution

Partition input into
random subsets



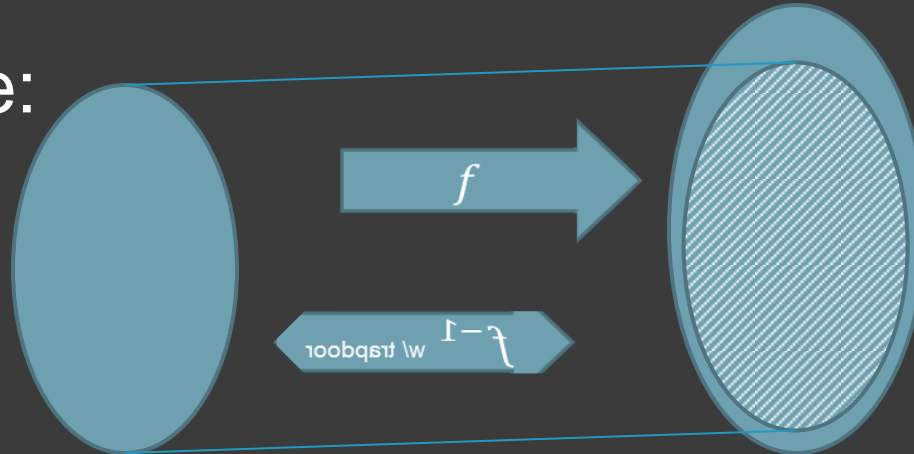
PRIV-IND \Rightarrow PRIV1-IND with Incrementality

Standard Model

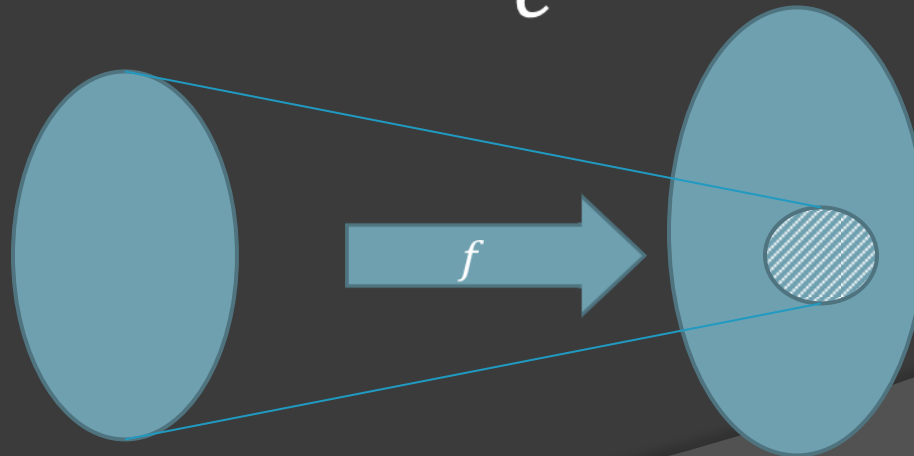
DDH \Rightarrow PRIV1-IND with Incrementality

Lossy Trapdoor Functions

Injective mode:



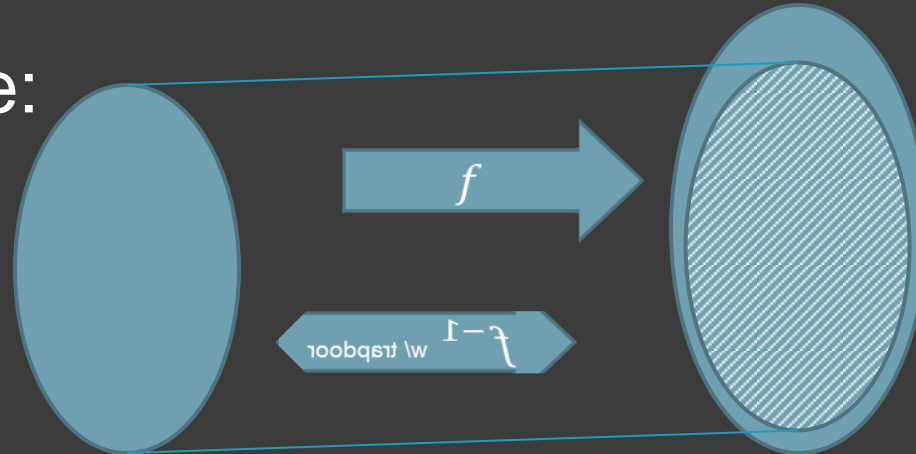
Lossy mode:



\approx_c

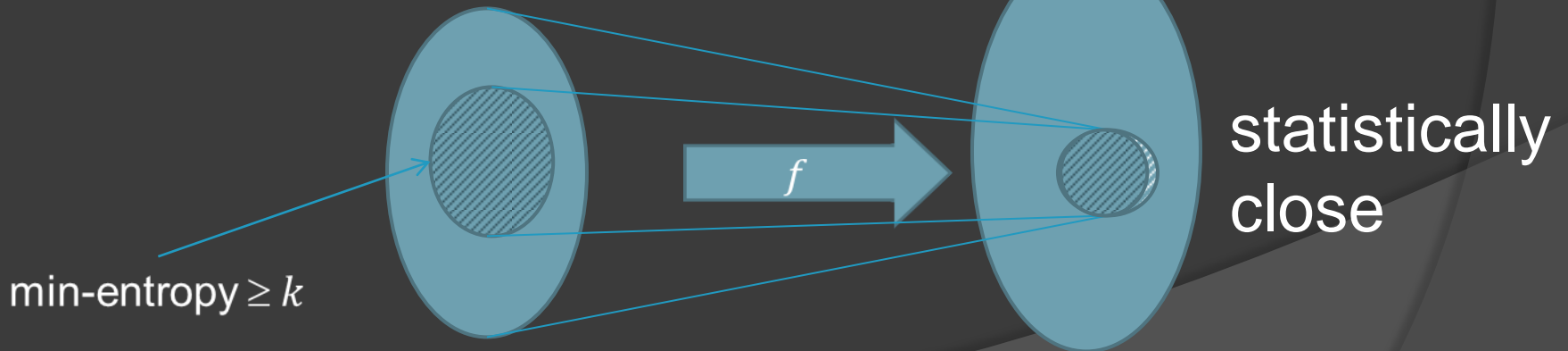
Smooth Trapdoor Functions

Injective mode:



\approx_c

Smooth mode:



Smooth Trapdoor Functions \Rightarrow PRIV1-IND

Security

$$\begin{array}{l} \text{injective mode:} \\ \text{smooth mode:} \end{array} \quad \begin{array}{ccc} \text{min-entropy} \geq k & & \text{min-entropy} \geq k \\ \downarrow & & \downarrow \\ f(M) & & f(N) \\ \approx_c & & \approx_c \\ f(M) & \approx & f(N) \end{array}$$

Construction of PRIV1-IND

Lossy Trapdoor Function

Pairwise-independent permutation

$$\odot \underbrace{f \circ \pi}$$

Smooth Trapdoor Function

Deterministic Public-Key Encryption

Construction of PRIV1-IND

Lossy Trapdoor Function

Pairwise-independent permutation



$$f \circ \pi$$

Smooth Trapdoor Function

Incremental Deterministic Public-Key Encryption

Construction of Lossy TDF

\mathbb{G} - group of order p generated by g

Key generation

- Sample $A \leftarrow \mathbb{Z}_p^{n \times n}$
- Output $sk = A^{-1}$ and $pk = g^A \in \mathbb{G}^{n \times n}$

$$(g^A)_{ij} = (g^{a_{ij}})$$

Encryption

- Given $\vec{m} \in \{0,1\}^n$ output $g^{A\vec{m}} \in \mathbb{G}^n$

Decryption

- Given $g^{\vec{v}} \in \mathbb{G}^n$ compute $g^{\vec{m}} = g^{A^{-1}\vec{v}} \in \mathbb{G}^n$
- Output $\vec{m} \in \{0,1\}^n$

Security Argument: Lossy TDF

$$\begin{array}{ccc} \text{rank } n & & \text{rank } 1 \\ & \swarrow & \searrow \\ & g^A \approx_c g^B & \end{array}$$

$$g^{A\vec{m}} \text{ — injective} \quad g^{B\vec{m}} \text{ — } \log p \text{ bits}$$

Towards Incremental Smooth TDF

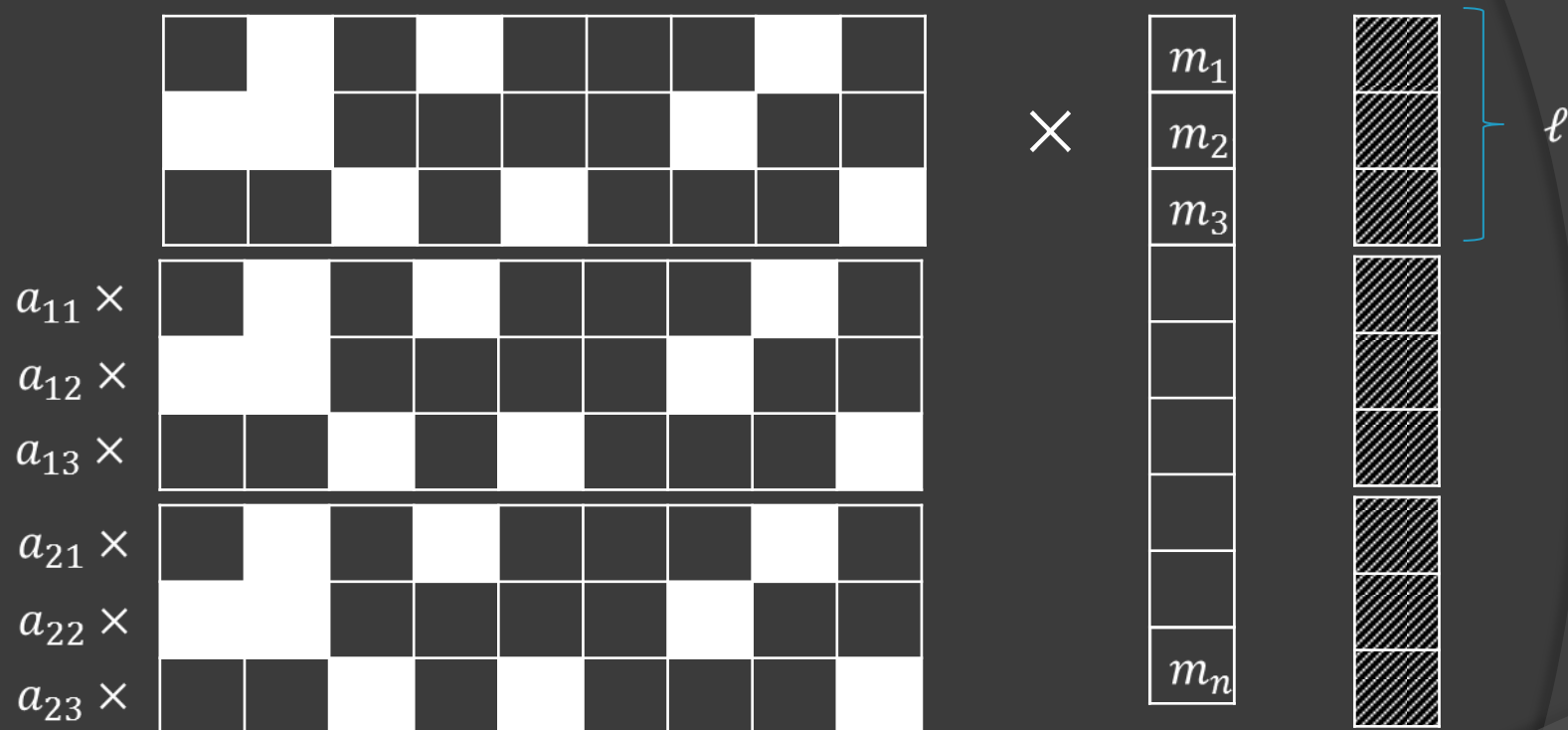
rank n
sparse

rank ℓ
sparse

$$g^A \approx_c g^B$$

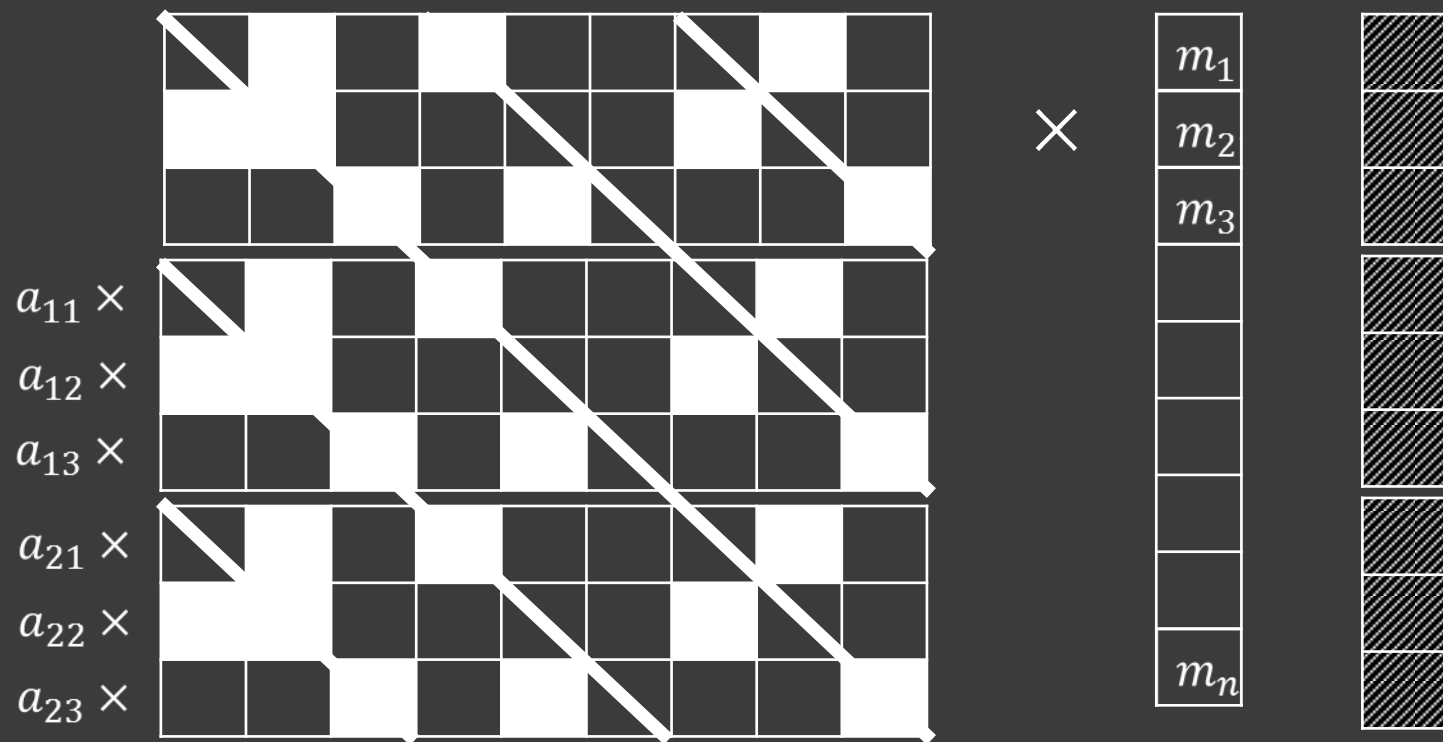
$g^{A\vec{m}}$ — injective if \vec{m} has min-entropy k ,
 $g^{B\vec{m}}$ statistically close to
the uniform over its range

Towards Incremental Smooth TDF

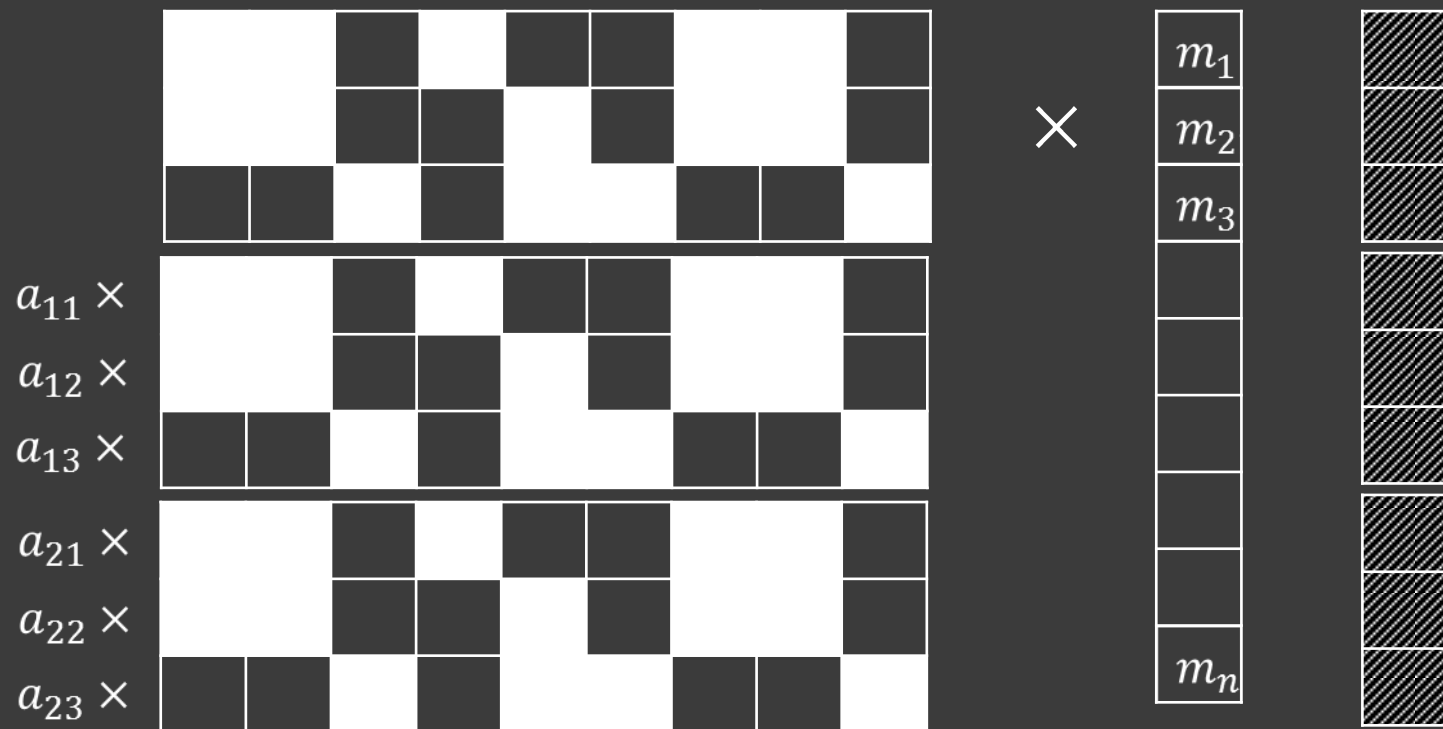


Sample-then-extract + Leftover Hash Lemma

Towards Incremental Smooth TDF

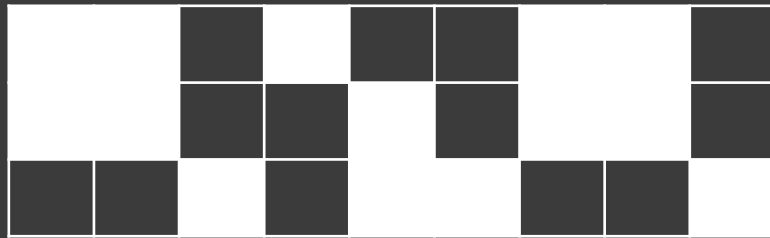


Towards Incremental Smooth TDF

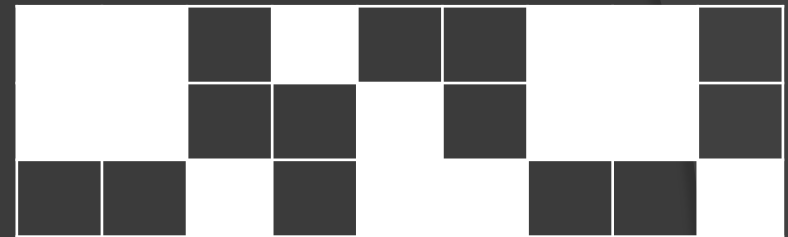


Smooth vs Injective Mode

rank ℓ



full rank



$a_{11} \times$

$a_{12} \times$

$a_{13} \times$

$a_{21} \times$

$a_{22} \times$

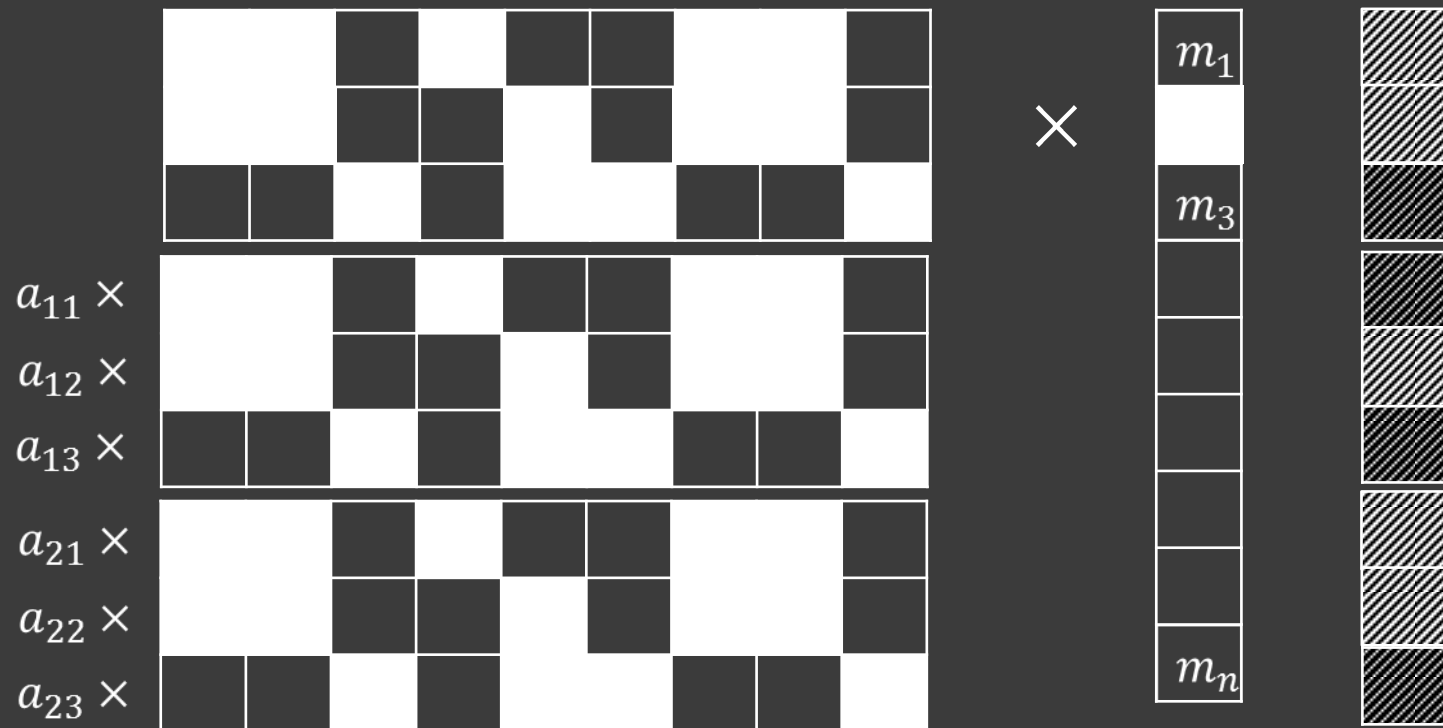
$a_{23} \times$

g

\approx_c

g

Incrementality



Open Problems

Incremental Deterministic Encryption:

- Stronger security: PRIV-IND (multiple messages)
- Length-preserving in the standard model

Deterministic Encryption:

- Relaxing the definition to allow dependency on the public key